### Views and View Management

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 $\begin{array}{c} {\rm CS~640} \\ {\rm Principles~of~Database~Management~and~Use} \\ {\rm Winter~2013} \end{array}$ 

These slides are based on slide sets provided by M. T. Öszu and by R. Ramakrishnan and J. Gehrke

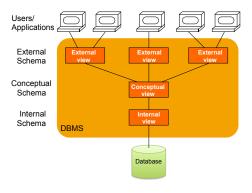
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### Recall the Three Levels of Abstraction



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### Views

### Definition (View)

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A *view* is a relation in the external schema whose instance is determined by instances of relations in the conceptual schema.

A view has many of the same properties as a base relation in the conceptual schema:

- its schema information appears in the database schema
- access controls can be applied to it
- other views can be defined in terms of it

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## Types of Views Notes Virtual view: Only the definition is stored • evaluating statements concerning a virtual view means rewriting those statements based on the view definition Materialized view: Instance is stored in the DBS • i.e. corresponding objects exist on the physical level CS 640 Views Winter 2013 4 / 15 Creating and Deleting Views using SQL Notes • General form: create [materialized] view <name> as <query> drop view <name> • Example: create view ResearchProjects as ( **select** projno, projname, lastname **as** empname from project, employee where respemp = empid and type = 'Research' ) CS 640 Views Winter 2013 5 / 15 Querying Views Notes Query a view as if it were a base relation. select projname from ResearchProjects where empname = 'Smith' What happens when you query a virtual view?

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### Recall: Query Processing Steps (View Resolution)

```
Parsing

Query Validation

View Resolution

View Resolution

View definitions

Optimization

Plan Compilation

Execution

As given in:

Goetz Graefe: Query Evaluation Techniques for Large
Databases. ACM Comp. Surveys 25(2): 73-170 (1993).
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Views

# Notes

### Querying Views (cont'd)

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# Notes

## Updating a Database via Views

- Modifications to a view's instance must be propagated back to instances of relations in conceptual schema.
- Some views cannot be updated unambiguously.
   Conceptual Schema

Ea	Can ad	ıan		
Dave	Canad	ian		
Wes	Ameri	can		
National	.Pastime	S		
CITIZENSHIP		PASTIME		
Canadian		Hockey		
Canadian		Curling		
American		Hockey		
American		Baseball		

NAME | CITIZENSHIP

Persons



### External Schema

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PersonalPastimes					
NAME	PASTIME				
Ed	Ho cke y				
Ed	Curling				
Dave	Ho cke y				
Dave	Curling				
Wes	Ho cke y				
Wes	Baseball				

- ① What does it mean to insert (Darryl, Hockey)?
- What does it mean to delete (Dave, Curling)?

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### View Updates in SQL

According to SQL-92, a view is updatable only if its definition satisfies a variety of conditions:

- The query references exactly one table
- The query only outputs simple attributes (no expressions)
- There is no grouping/aggregation/distinct
- There are no nested queries
- There are no set operations

These rules are more restrictive than necessary.

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### Motivation for Materializing Views

• Recall our example query:

select projname  $\textbf{from} \ \texttt{ResearchProjects}$ where empname = 'Smith'

• and the rewritten version:

select projname from project, employee where respemp = eid and type = 'Research' and lastname = 'Smith'

- Suppose we frequently ask this type of query, varying employees.
- Then, materializing the view might be a good idea.
- Additionally creating a clustered index on (the view's) attribute empname might be even better.

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### Motivation for Materializing Views (cont'd)

### Advantage

Using materialized views may improve query execution performance significantly (in particular for views with a complex definition).

### Downside

Need to maintain the view as the underlying base relations change.

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### Issues Related to Using Materialized Views

View design: What views should we materialize, and what indexes should we build on those views?

View exploitation: Given a query and a set of materialized views, can we use the materialized views to answer (parts of) the query more efficiently?

> • Different uses may be possible (another dimension for the optimizer)

View maintenance: How and when to refresh the materialized views?

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### Maintenance Policies

Immediate: Do the refresh as part of the transaction that modifies the underlying base relations.

- Pro: Materialized view is always consistent
- Con: Updates are slowed

Deferred: Do the refresh some time later, in a separate transaction.

- Pro: Can scale to maintain many views without slowing updates
- · Con: View may be out of date

Lazy Delay refresh until next query on view; then refresh before answering the query.

Periodic Refresh periodically. Widely used, especially for asynchronous replication in distributed databases, and for warehouse applications.

Event-based E.g., after a fixed number of updates to underlying base relations.

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### Summary and Outlook

- · Views are relations in the external schema whose instances are determined by relations in the conceptual schema
- · Types: virtual views, materialized views
- Updating a database via views may be ambiguous
  - First topic for our discussion next week A. M. Keller: The Role of Semantics in Translating View Updates. IEEE Computer 19(1): 63-73 (1986).
- · Materialized views may speed up query execution significantly
- Issues related to using materialized views:
  - View design
  - View maintenance
  - View exploitation ← second topic next week J.Goldstein and P.Larson: Optimizing Queries Using Materialized Views: A Practical, Scalable Solution. SIGMOD, 2001.

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